

1.204 Lecture 23

**Analytic approximations
Vehicle routing
Transit design**

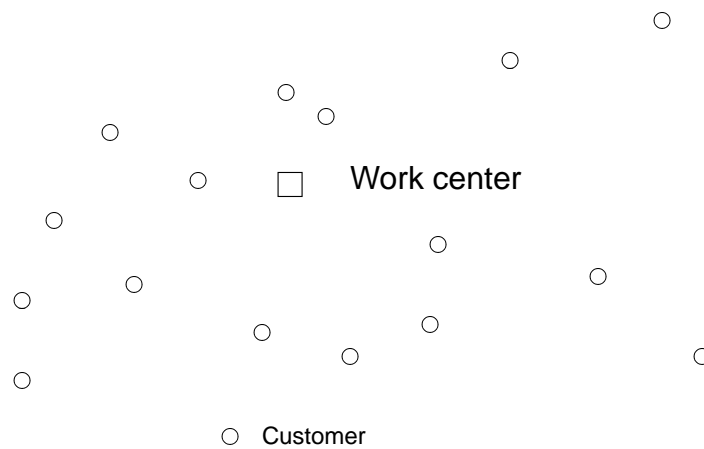
Analytic approximations

- **First spiral in developing problem solution**
 - Assist in requirements, prototyping, initial results, review
 - Many analytic approximations are visual, unlike almost all algorithms
 - Recall role of visualization in finding roots of equations, and how poorly algorithms do without it
 - Generally allow a broader treatment of the question, with more variables, more flexible objectives and constraints
 - Provide guidance in framing heuristics
 - Many real problems do not have optimal algorithms
 - We have very few $O(n)$ or $O(n^2)$ algorithms for complex problems; most are $O(2^n)$

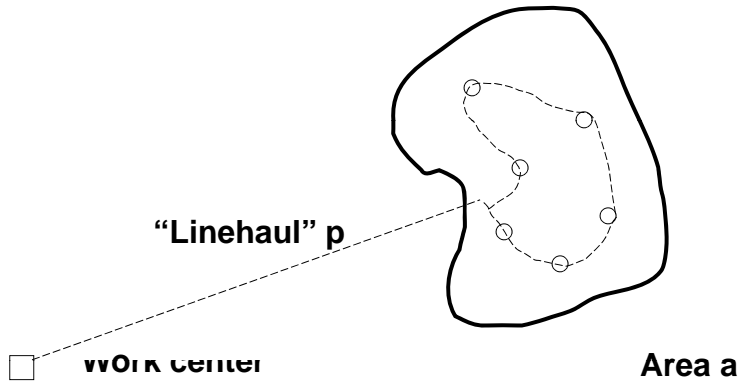
Vehicle routing

- **Variables:**
 - Number of routes or employees
 - Number of customers
 - Time windows or appointments
 - Capacity of vehicle or employee
 - Whether customers are known at start of route
 - And many others...
- **Objectives**
 - Customer service (timeliness, appointments)
 - Cost minimization
- **Constraints**
 - Labor rules, ...

Dispatch routing options



How to serve customers?

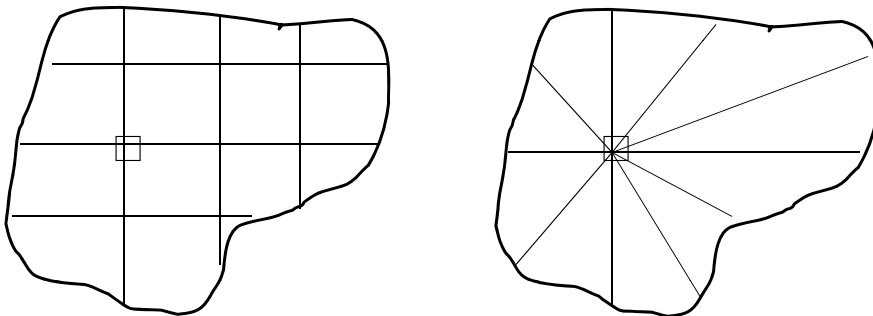


Trip length L to visit n randomly distributed customers in area a and return:

$$L = 2p + k \sqrt{na}$$

(Beardwood, Halton, Hammersley 1959)

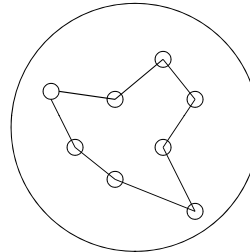
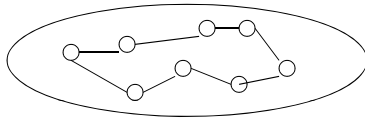
Shape of dispatch zones



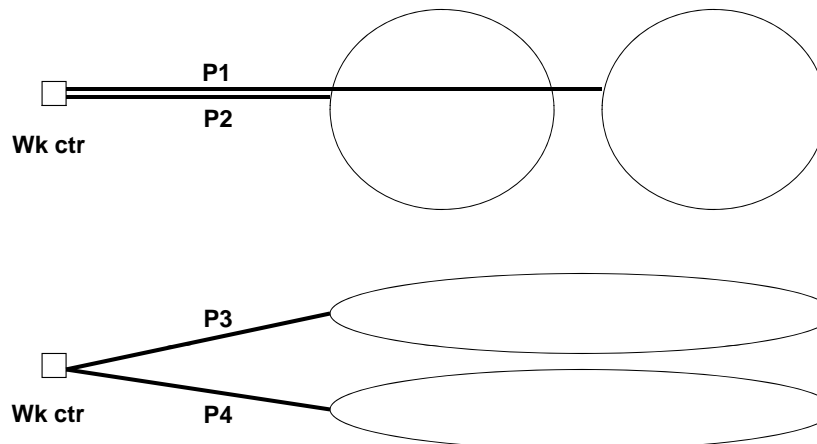
Which shape is better?

Zone shape

- **Let's try to elongate them**
 - Tour length is the same for same number of points, same area, different shapes

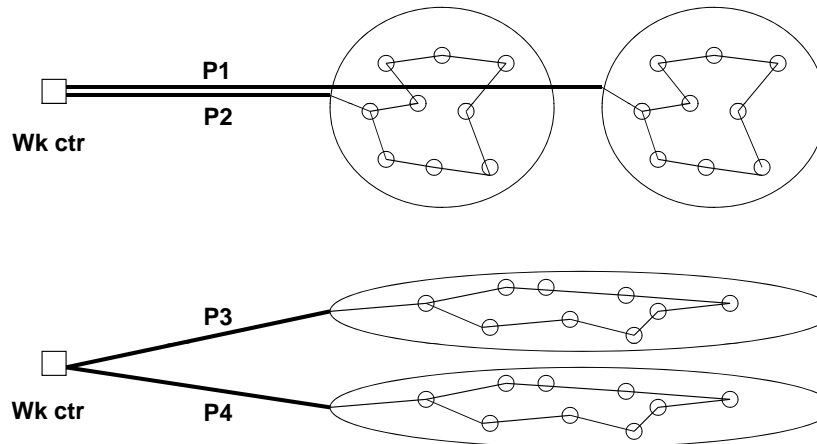


Comparison



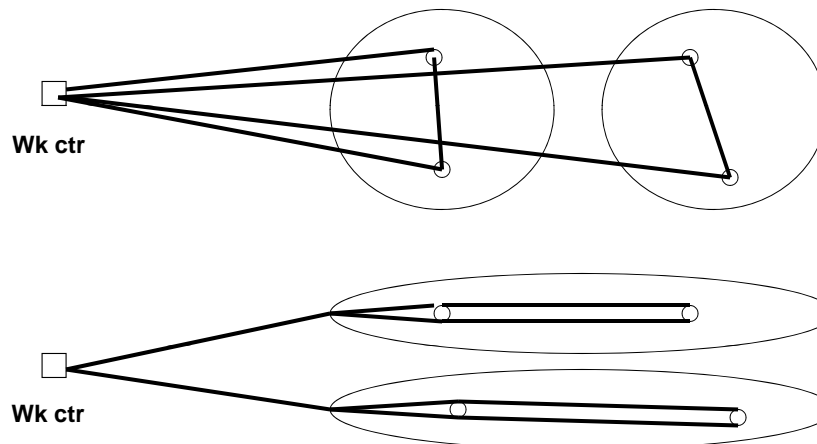
Elongated zones have shorter driving distance if there are a lot of customers in each zone

Elongated zones, many customers



With many customers, elongated zones are better

Fat zones, few customers



**If there are only a few customers, shape matters more
In this case, 'fat' zones are better**

Rules for building tours

- The “break point” for fat versus skinny zones is about 6 customers, based on simulation and geometric probability
 - If 6 or more customers can be served on a route:
 - Break up the area into skinny zones with the target number of points (6 or more)
 - Build tours in each zone, and fine tune
 - If 5 or fewer customers can be served on a route:
 - Break up the area into fat zones with the target number of points (5 or fewer). Only a few zones will touch the work center
 - Build tours in each zone, and fine tune
- Many dispatch systems ‘cluster’ jobs, which implicitly creates fat zones rather than skinny
 - Rack servicing has ~30-40 stops per day. Use skinny zones
 - Telecom dispatch has ~2-4 jobs per day. Use fat zones
 - Shared taxi has 2-4 stops per tour. Use fat zones
 - Dial-a-ride hopes to have 8-10 stops per tour. Use skinny zones

Rules for building tours-time windows

- Build skeleton elongated or fat routes (implicit zones) based on expected customer demand
 - Non-intersecting, non-overlapping routes
- Schedule stops in the following priority:
 - Tight time windows far from work center first
 - Then tight time windows near work center
 - Then other jobs far from work center
 - Then other jobs near work center
 - Pull next day’s work into today’s routes as feasible
 - Don’t give successive jobs with tight time windows to the same tech, if it can be avoided
- Rules determined from analysis and simulation

Summary- dispatch analysis

- **Done before writing dispatch algorithm or system**
 - Understand the problem, objectives and constraints
 - Use analytical optimization, simulation, probability, ...
 - Deal with broader set of issues than a single algorithm
 - Develop guidance for heuristics to be used

Transit system design

- **Variables for bus system design:**
 - Number of routes (route spacing)
 - Headway (frequency of service)
 - Fare
 - Vehicle size
 - Route length
 - Bus stop spacing
 - Express versus local service
 - Transfer pattern

Transit system design

- **Objectives:**
 - Maximize ridership
 - Minimize deficit (or maximize profit)
 - Equity in service levels
- **Constraints**
 - Available resources (deficit limit)
 - Minimum service levels
 - System capacity

Route spacing and headway

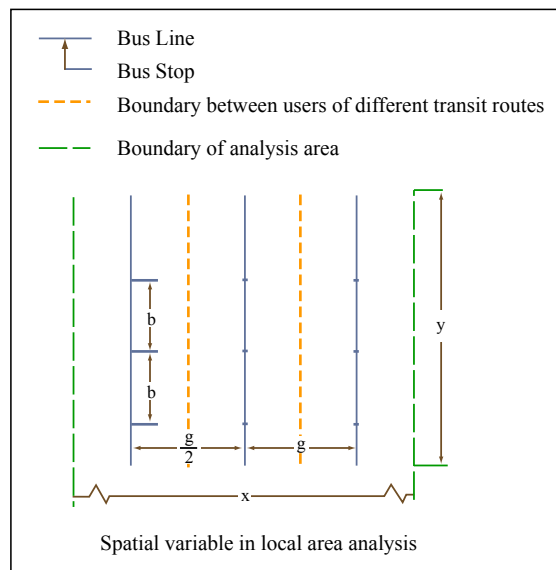
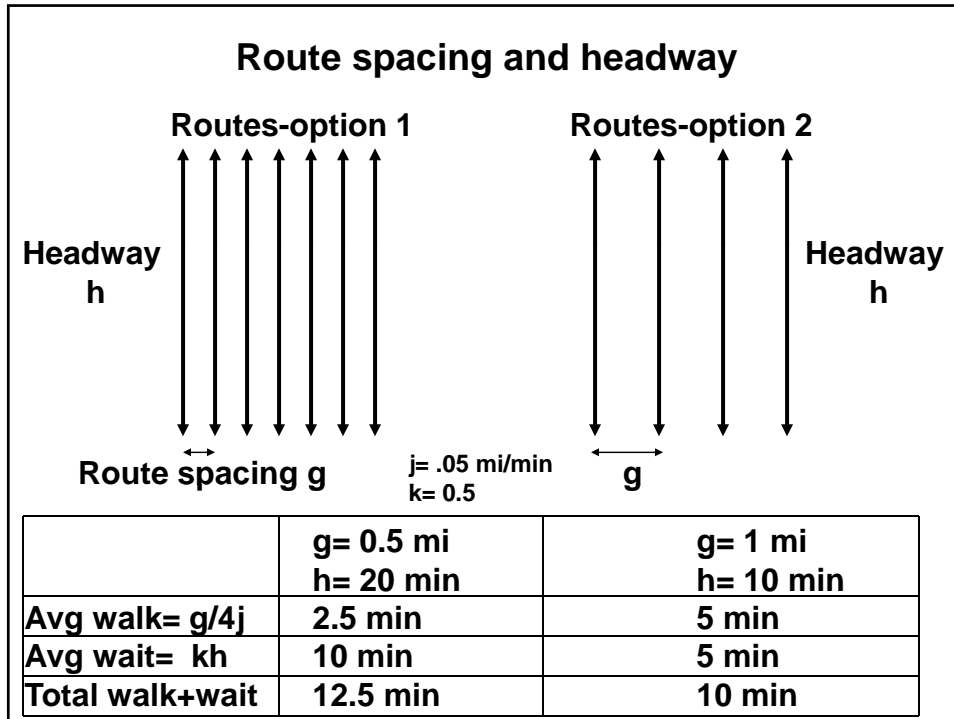


Figure by MIT OpenCourseWare.



- ### Route spacing and headway
- **General result over many objectives and constraints:**
 - **Optimal route spacing and headway are related by**
 - $h^* = g / 4jk$
 - **At this point, average walk time= average wait time**
 - **Complications:**
 - Ratio of wait time/headway, k , may vary with headway
 - There may be a 'walk refusal distance', and demand response to walk distance may be nonlinear
 - Headway varies over the day on a route: either choose an average spacing, or have peak-only routes
 - **Option 2 on previous slide is 'optimal'**
 - Same operating cost and capacity as option 1
 - Better service (lower sum of walk and wait times)
 - Higher ridership

Fare and demand function

- Introduce a demand function:

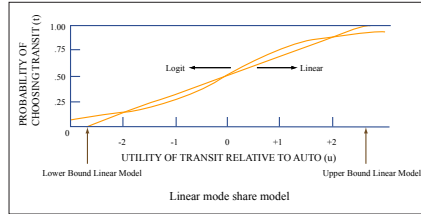


Figure by MIT OpenCourseWare.

- Mode share $t = (a_1 + a_2(kh + (g+b)/4)) + a_3 * d/v + a_4 * f + a_5 * d$
 - Where
 - $a_1..a_5$ are demand coefficients (a_5 is auto coefficient)
 - d is route distance, v is bus velocity, f is bus fare
- Total bus ridership in area $P = TpXYt$
 - Where
 - X, Y are dimensions (mi),
 - p is trip density (trips/mi²/min),
 - T is time period (min)

Cost and objective function

- Bus operating costs:
 - $C = 2 XYTc/ghv$
 - There are X/g routes operating $2T/h$ trips of length Y/v at unit cost per minute of c
- Objective function: maximize net social benefits
 - Max consumers' surplus G + revenue $R (=Pf)$ – cost C
 - Subject to a deficit constraint ($C - R \leq M$)
 - G is a proxy for external benefits (air quality, GHG, ...)

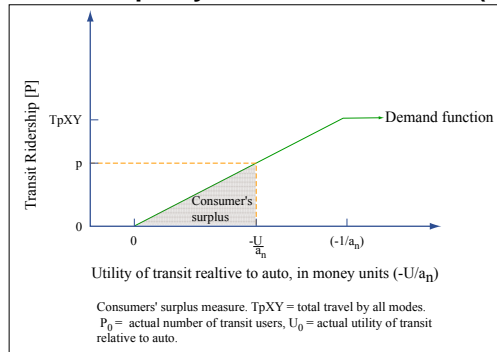


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Overall formulation

- Max net benefit subject to (cost – revenue) <= M

$$\max B_2 = -TpXY(a_1 + a_2(kh + (g + b)/4j) + a_3d/v + a_4f + a_5d)^2/2a_4$$

subject to

$$2XTcY/ghv - TpXYf(a_1 + a_2(kh + (g + b)/4j) + a_3d/v + a_4f + a_5d) - M \leq 0, \quad g, h, f \geq 0.$$

- Formulate using Lagrange multiplier

$$\begin{aligned} \max B_1 = & -TpXY(a_1 + a_2(kh + (g + b)/4j) \\ & + a_3d/v + a_4f + a_5d)^2/2a_4 - y_1[2XTcY/ghv \\ & - TpXYf(a_1 + a_2(kh + (g + b)/4j) \\ & + a_3d/v + a_4f + a_5d)], \quad g, h, f \geq 0. \end{aligned}$$

Solutions

- Take derivatives with f, g, h and y2 to obtain 4 nonlinear equations in 4 unknowns, and solve approximately:

$$g^* \approx (32j^2ka_4c(2y_2 - 1)/vpa_2Ay_2)^{1/3}$$

$$h^* \approx (a_4c(2y_2 - 1)/2jk^2vpa_2y_2A)^{1/3}$$

$$f^* \approx [(1 - y_2)/(2y_2 - 1)]$$

$$\cdot [A/a_4 + (4kca_2^2(2y_2 - 1)/vpja_4^2Ay_2)^{1/3}].$$

- where
 - f is fare, g is route spacing, h is headway,
 - y2 is Lagrange multiplier or shadow price of \$1 of benefit relative to \$1 of deficit
- We vary y2 to get solutions ranging from min deficit (y2 infinite) to max social benefit (y2= 1)
 - We can also add a vehicle capacity constraint (y3)

Model summary

- **Model implemented in Java code**
 - Download code and documentation
- **Provides framework for designing bus system:**
 - Routes, headways, fares, vehicle sizes, express/local service
 - Bus stop spacing (fewer are better)
 - Route circuitry (less circuitry is better)
 - (Model variation used in planning Logan Express)
- **Allows variation in objective and constraints**
- **Provides insight before addressing detailed system design with actual network and routes, using optimization algorithms and simulation**
 - Most of the the term was spent on optimization algorithms for decisions and design
 - Simulation not covered, used for truly difficult/detailed issues
- **We'll do analytical approximations for queuing systems in the next lecture**

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1.204 Computer Algorithms in Systems Engineering
Spring 2010

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