

## 1.204 Computer Algorithms in Systems Engineering

Spring 2010

### Problem Set 4: Satellite data sets

**Due: 12 noon, Monday March 29, 2010**

Please see the code provided for the implementation. There are three separate solutions in the Java zip file. All use identical versions of Merge, SatelliteData and MinHeap.

- Solution with an explicit tree uses MergeMain and TreeNode
- Solution without an explicit tree uses MergeMainNoTree and Node
- Solution used to answer question 3 below uses MergeMainNoTreeAnalysis and Node

Analysis questions:

- In a full binary tree, each node has either zero or two children. In each step of the algorithm, we add two children (which are either nodes or trees) to a node or a tree. In the initial step, we add two nodes to a single node and thus creating a full binary tree. Adding two nodes to a full binary tree also creates another full binary tree. Joining two full binary trees also creates another full binary tree. Thus, at the end of the optimal merge algorithm, we have a full binary tree.
- For a binary heap, the insert operation takes  $O(\lg n)$  steps, and the deletion operation also takes  $O(\lg n)$  steps. Creating the initial heap of all the nodes takes  $O(n)$  steps, since it uses heapify. Thus the algorithm takes:
  - $O(n)$  steps to create the heap initially
  - Then in each step of the algorithm, we have two delete operations and one insert operation on the heap. After each step, we reduce the number of elements in the heap by 1.
  - Thus the total run time is:

$$O(n) + \sum [O(\lg(n-1)) + O(\lg(n-2)) + O(\lg(n-1))]$$

The second term in the above expression is  $O(n \lg n)$ , since

$$\lg(n) + \lg(n-1) + \dots = \lg(n!),$$

and by Stirling's formula,

$$\lg(n!) \sim n \lg(n) + cn, \text{ where } c \text{ is a constant.}$$

Thus the run time for this algorithm is  $O(n \lg n)$ .

3. See code provided. An example run is shown below. Note that the number of steps per data element increases essentially exactly as  $\lg n$ , which means the algorithm performance is indeed almost exactly  $O(n \lg n)$ , where the number of data elements is  $n$ . In this example run, each dataset has exactly 200 randomly generated elements.

Arrays	Steps	Data	Ratio
1	0	200	0.0
2	400	400	1.0
3	1000	600	1.6666666
4	1600	800	2.0
5	2400	1000	2.4
6	3200	1200	2.6666667
7	4000	1400	2.857143
8	4800	1600	3.0
9	5800	1800	3.2222223
10	6800	2000	3.4
11	7800	2200	3.5454545
12	8800	2400	3.6666667
13	9800	2600	3.7692308
14	10800	2800	3.857143
15	11800	3000	3.9333334
16	12800	3200	4.0
17	14000	3400	4.117647
18	15200	3600	4.2222223
19	16400	3800	4.3157897
20	17600	4000	4.4

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